Title: Solving Problems by Searching
AIMA: Chapter 3 (Sections 3.4, 3.5, and 3.6)

Introduction to Artificial Intelligence
CSCE 476-876, Spring 2004
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function GENERAL-SEARCH(*problem, strategy*) **returns** a solution, or failure
  initialize the search tree using the initial state of *problem*
  loop do
    if there are no candidates for expansion **then return** failure
    choose a leaf node for expansion according to *strategy*
    if the node contains a goal state **then return** the corresponding solution
    else expand the node and add the resulting nodes to the search tree
  end

Essence of search: which node to expand first?
$\longrightarrow$ search strategy

A strategy is defined by picking the *order of node expansion*

Types of Search

Uninformed: use only information available in problem definition

Heuristic: exploits some knowledge of the domain

Uninformed search strategies

1. Breadth-first search
2. Uniform-cost search
3. Depth-first search
4. Depth-limited search
5. Iterative deepening depth-first search
6. Bidirectional search

Search strategies

Criteria for evaluating search:

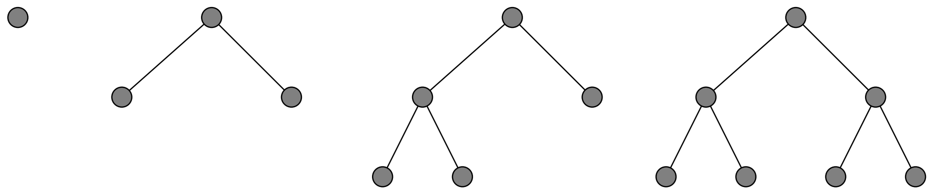
1. Completeness: does it always find a solution if one exists?
2. Time complexity: number of nodes generated/expanded
3. Space complexity: maximum number of nodes in memory
4. Optimality: does it always find a least-cost solution?

Time/space complexity measured in terms of:

- b: maximum branching factor of the search tree
- d: depth of the least-cost solution
- m: maximum depth of the search space (may be ∞)

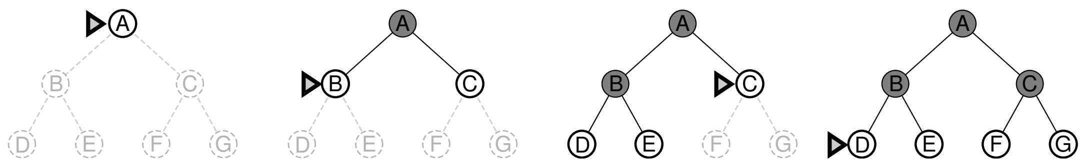
Breadth-first search (I)

→ Expand root node
→ Expand *all* children of root
→ Expand *each* child of root
→ Expand successors of each child of root, etc.



⟶ Expands nodes at depth d before nodes at depth $d + 1$
⟶ Systematically considers all paths length 1, then length 2, etc.
⟶ Implement: put successors at end of queue.. FIFO

Breadth-first search (2)



Breadth-first search (3)

$\longrightarrow$ One solution?

$\longrightarrow$ Many solutions? Finds shallowest goal first

1. Complete? Yes, if b is finite
2. Optimal? provided cost increases monotonically with depth, not in general
3. Time? $1 + b + b^2 + b^3 + \ldots + b^d + b(b^d - 1) = O(b^{d+1})$

   $$O(b^{d+1}) \begin{cases} \text{branching factor } b \\ \text{depth } d \end{cases}$$
4. Space? same, $O(b^{d+1})$, keeps every node in memory, big problem
   can easily generate nodes at 10MB/sec so 24hrs = 860GB

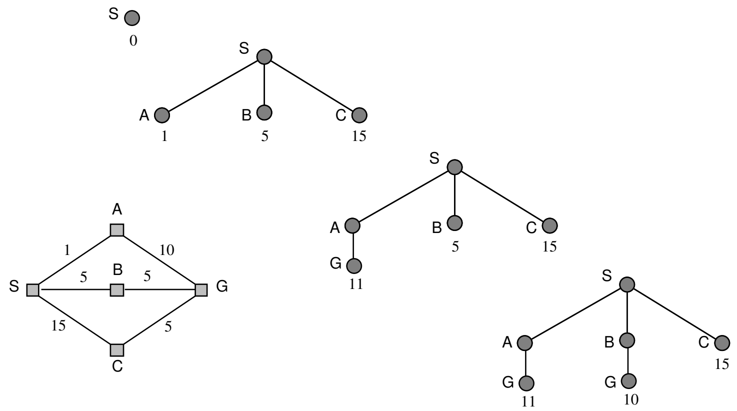
Uniform-cost search (I)

$\longrightarrow$ Breadth-first does not consider <u>path cost</u> $g(x)$

$\longrightarrow$ Uniform-cost expands first lowest-cost node on the fringe

$\longrightarrow$ Implement: sort queue in decreasing cost order

When $g(x) = \text{Depth}(x) \longrightarrow$ Breadth-first $\equiv$ Uniform-cost



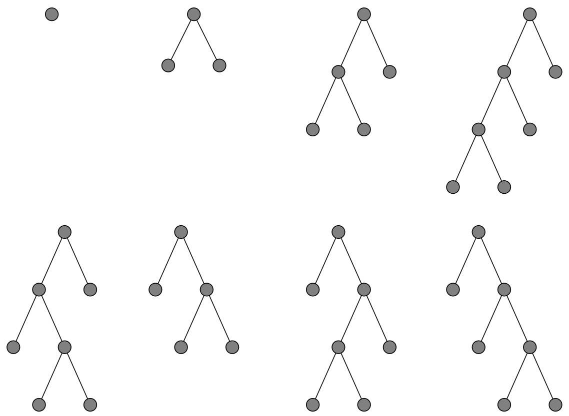
(a) (b)

Uniform-cost search (2)

1. Complete?
   Yes, if cost $\geq \epsilon$
2. Optimal?
   If the cost is a monotonically increasing function
   When cost is added up along path, an operator's cost ....?
3. Time?
   # of nodes with $g \; \leq$ cost of optimal solution, $O(b^{\lceil C^*/\epsilon \rceil})$
   where C^* is the cost of the optimal solution
4. Space?
   # of nodes with $g \; \leq$ cost of optimal solution, $O(b^{\lceil C^*/\epsilon \rceil})$

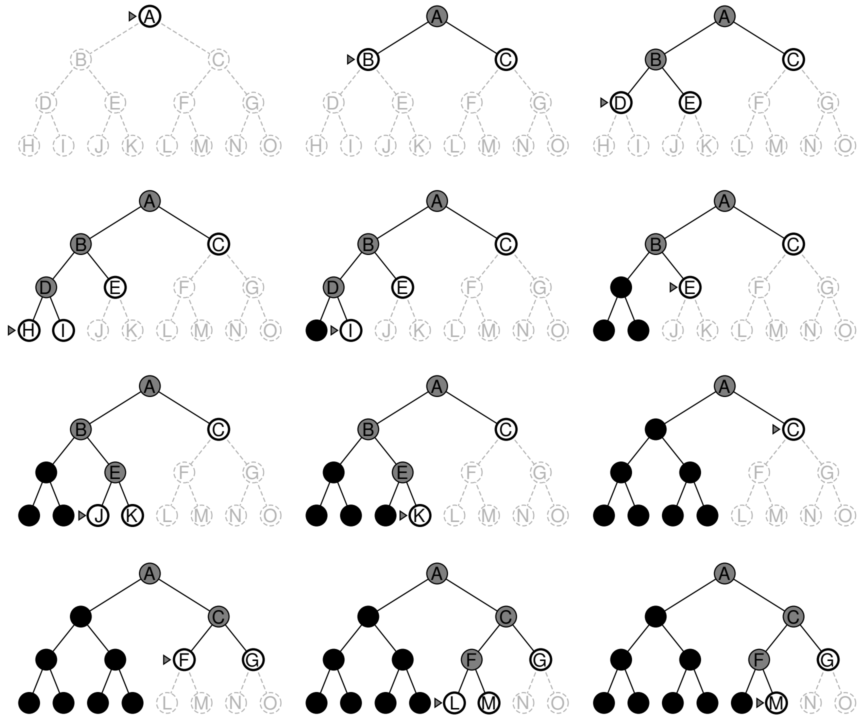
Depth-first search (I)

⟶ Expands nodes at deepest level in tree
⟶ When dead-end, goes back to shallower levels
⟶ Implement: put successors at front of queue.. LIFO



⟶ Little memory: path and unexpanded nodes
For b: branching factor, m: maximum depth, space ....?

Depth-first search (2)



Depth-first search (3)

Time complexity:

We may need to expand all paths, $O(b^m)$

When there are many solutions, DFS may be quicker than BFS

When m is big, much larger than d, ∞ (deep, loops), .. troubles

$\longrightarrow$ Major drawback of DFS: going deep where there is no solution..

Properties:

1. Complete? No in infinite-spaces, complete in finite spaces
2. Optimal?
3. Time? $O(b^m)$ Woow..
   terrible if m is much larger than d, but if solutions are dense, may be much faster than breadth-first
4. Space? $O(bm)$, linear! Woow..

Depth-limited search (I)

⟶ DFS is going too deep, put a threshold on depth!
For instance, 20 cities on map for Romania, any node deeper than 19 is cycling. Don't expand deeper!

⟶ Implement: nodes at depth l have no successor

Properties:

1. Complete?
2. Optimal?
3. Time? (given l depth limit)
4. Space? (given l depth limit)

Problem: how to choose l?

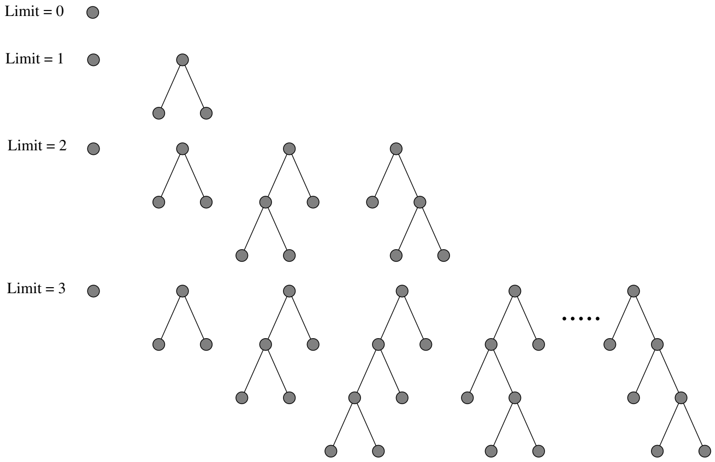
Iterative-deepening search (I)

$\rightarrow$ DLS with depth $= 0$

$\rightarrow$ DLS with depth $= 1$

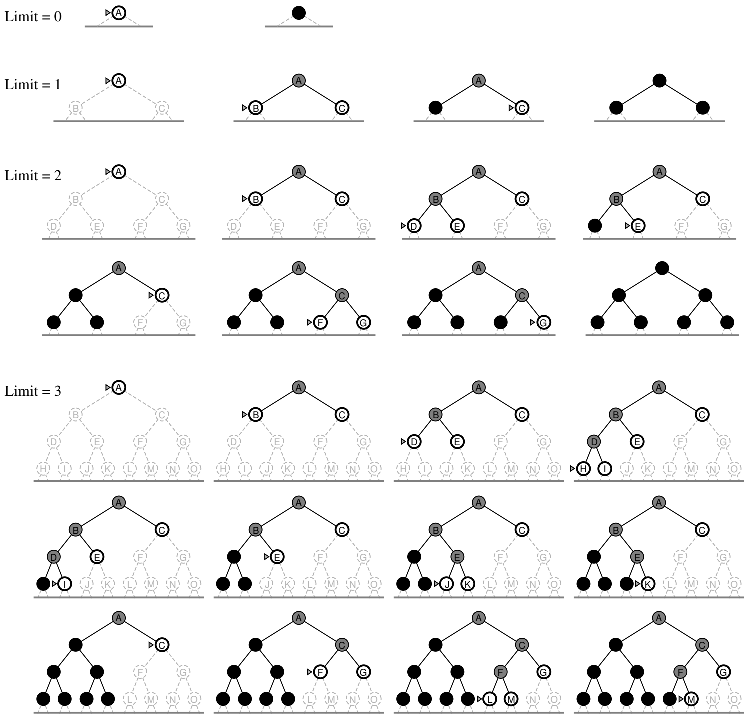
$\rightarrow$ DLS with depth $= 2$

$\rightarrow$ DLS with depth $= 3$...



$\longrightarrow$ Combines benefits of DFS and BFS

Iterative-deepening search (2)



Iterative-deepening search (3)

$\longrightarrow$ combines benefits of DFS and BFS

Properties:

1. Time? $(d+1).b^0 + (d).b + (d-1).b^2 + \ldots + 1.b^d = O(b^d)$
2. Space? $O(bd)$, like DFS
3. Complete? like BFS
4. Optimal? like BFS (if step cost = 1)

Iterative-deepening search (4)

$\longrightarrow$ Some nodes are expanded several times, wasteful?

N(BFS) $= b + b^2 + b^3 + \ldots + b^d + (b^{d+1} - d)$

N(IDS) $= (d)b + (d-1)b^2 + \ldots + (1)b^d$

Numerical comparison for $b = 10$ and $d = 5$:

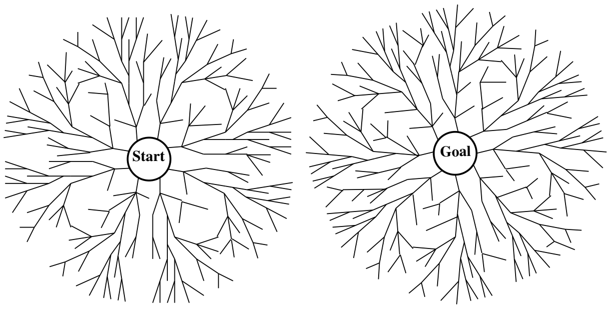
N(IDS) = 50 + 400 + 3,000 + 20,000 + 100,000 = 123,450

N(BFS) = 10 + 100 + 1,000 + 10,000 + 100,000 + 999,990 = 1,111,100

$\longrightarrow$ IDS is preferred when search space is large and depth unknown

Bidirectional search (I)

→ Given initial state and the goal state, start search from both ends and meet in the middle



→ Assume same b branching factor, $\exists$ solution at depth d, time: $O(2b^{d/2}) = O(b^{d/2})$

$b = 10, d = 6$, DFS= 1,111,111 nodes, BDS=2,222 nodes!

Bidirectional search (2)

In practice :—(

- Need to define predecessor operators to search backwards
  If operator are invertible, no problem
- What if $\exists$ many goals (set state)?
  do as for multiple-state search
- need to check the 2 fringes to see how they match
  need to check whether any node in one space appears in the other space (use hashing)
  need to keep all nodes in a half in memory $O(b^{d/2})$
- What kind of search in each half space?

Summary

Criterion	Breadth-First	Uniform-Cost	Depth-First	Depth-Limited	Iterative Deepening
Complete?	Yes*	Yes*	No	Yes, if $l \geq d$	Yes
Time	b^{d+1}	$b^{\lceil C^*/\epsilon \rceil}$	b^m	b^l	b^d
Space	b^{d+1}	$b^{\lceil C^*/\epsilon \rceil}$	bm	bl	bd
Optimal?	Yes*	Yes*	No	No	Yes

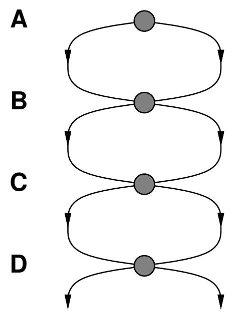
b branching factor
d solution depth
m maximum depth of tree
l depth limit

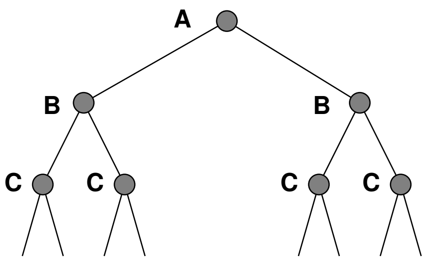
Loops: Avoid repeated states (I)

Avoid expanding states that have already been visited

Valid for both infinite and finite trees

$$\text{Example: } \begin{cases} m \text{ maximum depth} \\ m+1 \text{ states} \\ 2^m \text{ possible branches (paths)} \end{cases}$$



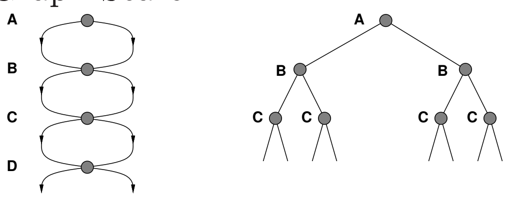


Loops: (2)

Keep nodes in two lists: $\begin{cases}\text{Open list: Fringe}\\ \text{Closed list: Leaf and expansed nodes}\end{cases}$

Discard a current node that matches a node in the closed list

Tree-Search $\longrightarrow$ Graph-Search



Issues:

1. Implementation: hash table, access is constant time
   Trade-off cost of storing+checking vs. cost of searching
2. Losing optimality
   when new path is cheaper/shorter of the one stored
3. BFS and IDS now require exponential storage

Summary

Path: sequence of actions leading from one state to another

Partial solution: a path from an initial state to another state

Search: develop a sets of partial solutions

- Search tree & its components (node, root, leaves, fringe)
- Data structure for a search node
- Search space vs. state space
- Node expansion, queue order
- Search types: uninformed vs. heuristic
- 6 uninformed search strategies
- 4 criteria for evaluating & comparing search strategies

Searching with partial information (I)

So far, we assumed:

- Environment fully observable
- Environment deterministic
- Agent knows effects of actions

Thus, agent

- always knows where it is
- can compute state where it will be after a sequence of actions

What happens when knowledge about states and actions is incomplete?

Searching with partial information (2)

Incompleteness yields 3 types of problems:

- Sensorless (conformant) problems
- Contingency problems
- Exploration problems

Sensorless problems (conformant)

- Environment not observable, no percepts
- Agent does not know in which exact state it is
  - agent may be in one of more possible initial states
  - an action may lead to one or more possible successor states

Contingency problems

- environment partially observable or actions are uncertain
- agent's percepts provide new input after each action, a contingency to plan for
- **Adverserial problems:** uncertainty caused by action of other agents

Exploration problems

- States and actions of the environment are unknown
- Agent must act to discover them
- Extreme case of contingency problem

Sensorless problems (I)

Vacuum cleaner: no sensors, but agent knows effects of actions

Agent may be in any state $\{1, 2, 3, 4, 5, 6, 7, 8\}$

- $[Right]$ always ends in $\{2, 4, 6, 8\}$
- $[Right, Suck]$ always ends in $\{4, 8\}$
- $[Right, Suck, Left, Suck]$ always works, coerces the world into 7

Sensorless problems (2)

Environment not (fully) observable:

- Agent must think about sets of states,
- Agent has a belief state (set of possible states)

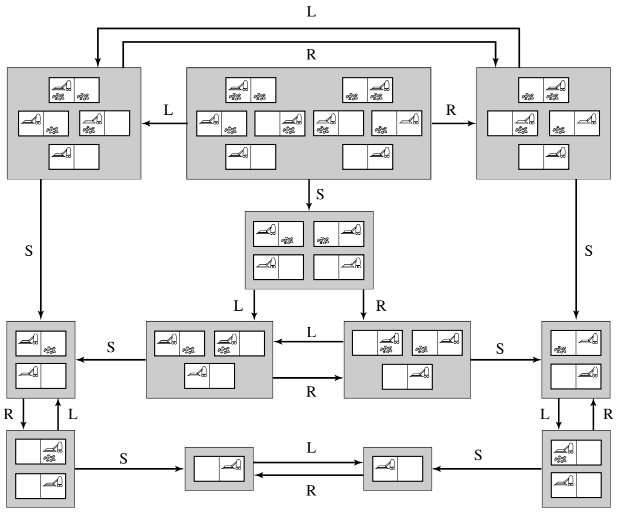
Environment fully observable: 1 belief state has 1 state

Solving sensorless problems: search in space of beliefs

- initial state is a belief state (all possible states)
- actions map 1 belief state into another
- belief state is union of applying action to each state in initial belief state
- goal is reached when all states in belief state are goal states

Sensorless problems (2)

vacuum cleaner: 12 belief states



In general:

8 states, 2^8 possible belief states

S states, 2^S possible belief states

Sensorless problems (3)

So far assumed deterministic environment
Approach/results hold for nondeterministic environment

Example: Murphy's law, *Suck* sometimes deposits dirt on carpet but <u>only if there is no dirt there already</u>

- [*Suck*] applied to State 4 leads to {2, 4}
- [*Suck*] applied to {1, 2, 3, 4, 5, 6, 7, 8} leads to ...
- Problem is unsolvable (Exercise 3.18)!!
  Agent cannot tell whether state is dirty and cannot predict whether *Suck* is going to make it dirty or clean

Contingency problems (I)

Environment partially observable or actions are uncertain

When agent can get some information:

- about environment
- from sensors
- after acting

Solution to a contingency problem is not a path, but a tree
$\longrightarrow$ branches are selected depending on percepts

Contingency problems (2)

Example: vacuum cleaner

- has 'local dirt' sensor, no 'remote dirt' sensor
- has location sensor
- Murphy's law

Now,

- Agent perceives $[L, Dirty]$, thinks in state {1, 3}
- Action $[Suck]$ leads to {5, 7}
- Action $[Suck, Right]$ leads to {6, 8}
- Action $[Suck, Right, Suck]$ leads to {8, 6}
  Plan can succeed (8), or fail (6)

Thus, action $[Suck, Right, \mathbf{if}[R, Dirty]\mathbf{then}Suck]$ leads to {8, 6}
Solution is a tree

Contingency problems (3)

Example: vacuum cleaner

- has 'local dirt' sensor and 'remote dirt' sensor
- has location sensor (fully observable)
- Murphy's law

Solution is a sequence of actions

Agent can proceed...

Contingency problems (4)

In general, agent

- acts before having a guaranteed plan (solution is a tree)
- needs to consider every possibility that might arise
  $\longrightarrow$ may be an overkill

It is (sometimes) necessary to start acting,
and deal with contingencies as they arise

- $\longrightarrow$ Interleave Search and Execution
- $\longrightarrow$ Useful for game playing and exploration problems