

Temporal Databases

(Chapter 5)

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Allen's Relations

Example Binary Relations between events I and J:

I Before *J* – when *I* is over before *J* starts.

I Contains *J* – when *J* occurs only when *I* occurs.

I Equals *J* – when *I* and *J* occur always simultaneously.

I Meets *J* – when *I* just finishes before *J* starts.

I Overlap *J* – when *I* and *J* overlap in time.

Axioms on Allen's Relations

Transitivity Axioms:

If I_1 Before I_2 and I_2 Before I_3 , then I_1 Before I_3 .

If I_1 Contains I_2 and I_2 Contains I_3 , then I_1 Contains I_3 .

If I_1 Equals I_2 and I_2 Equals I_3 , then I_1 Equals I_3 .

Meets-Before Axiom:

If I_1 Meets I_2 and I_2 Meets I_3 , then I_1 Before I_3 .

Representing Allen's Relations by Constraints

Let the beginning and the end of event I be I^- and I^+ Then we have:

| Relation of I and J | $I^- \quad J^-$ | $I^- \quad J^+$ | $I^+ \quad J^-$ | $I^+ \quad J^+$ |
|---------------------|-----------------|-----------------|-----------------|-----------------|
| After : | | $>$ | | |
| Before : | | | $<$ | |
| Meets : | | | $=$ | |
| Met by : | | $=$ | | |
| During : | $>$ | | | $<$ |
| Contains : | $<$ | | | $>$ |
| Equal : | $=$ | | | $=$ |
| Finishes : | $>$ | | | $=$ |
| Finished by : | $<$ | | | $=$ |
| Starts : | $=$ | | | $<$ |
| Started by : | $=$ | | | $>$ |
| Overlaps : | $<$ | | $>$ | $<$ |
| Overlapped by : | $>$ | $<$ | | $>$ |

Temporal Data Abstraction

View Level: Anderson at AT&T: 1980_____1997
Brown at IBM:_____1985_____1996
Clark at Lotus:_____1990__1991

Logical Level:

Employee

| Name | Company | Year |
|----------|---------|------|
| Anderson | AT&T | 1980 |
| ⋮ | ⋮ | ⋮ |
| Anderson | AT&T | 1997 |
| Brown | IBM | 1985 |
| ⋮ | ⋮ | ⋮ |
| Brown | IBM | 1996 |
| Clark | Lotus | 1990 |
| Clark | Lotus | 1991 |

Temporal Database Queries (Point-Based Queries)

Example: Find the employees who did not work during all the time that Brown worked.

```
SELECT      E1.Name  
FROM        Employee AS E1, Employee AS E2  
WHERE       E1.Name ≠ "Brown" AND  
            E2.Name = "Brown" AND  
            E2.Year NOT IN      (SELECT E3.Year  
                                  FROM   Employee AS E3  
                                  WHERE  E3.Name = E1.Name)
```

- Point-based queries use the logical level.
- Point-based queries are usually only *standard SQL queries*.

Temporal Data Abstraction

View Level: Anderson at AT&T: 1980_____1997
 Brown at IBM: _____1985_____1996
 Clark at Lotus: _____1990__1991

Logical Level:

Employee

| Name | Company | Year |
|----------|---------|------|
| Anderson | AT&T | 1980 |
| ⋮ | ⋮ | ⋮ |
| Anderson | AT&T | 1997 |
| Brown | IBM | 1985 |
| ⋮ | ⋮ | ⋮ |
| Brown | IBM | 1996 |
| Clark | Lotus | 1990 |
| Clark | Lotus | 1991 |

Constraint Level:

Option 1: Use *order constraints*.

Employee

| Name | Company | Year | |
|----------|---------|------|----------------------------|
| Anderson | AT&T | t | $1980 \leq t, t \leq 1993$ |
| Anderson | AT&T | t | $1994 \leq t, t \leq 1997$ |
| Brown | IBM | t | $1985 \leq t, t \leq 1996$ |
| Clark | Lotus | t | $1990 \leq t, t \leq 1991$ |

Option 2 (TQuel): Use *From* and *To*.

Employee

| Name | Company | From | To |
|----------|---------|------|------|
| Anderson | AT&T | 1980 | 1993 |
| Anderson | AT&T | 1994 | 1997 |
| Brown | IBM | 1985 | 1996 |
| Clark | Lotus | 1990 | 1991 |

Querying TQuel Temporal Databases

Employee

| Name | Company | From | To |
|----------|---------|------|------|
| Anderson | AT&T | 1980 | 1993 |
| Anderson | AT&T | 1994 | 1997 |
| Brown | IBM | 1985 | 1996 |
| Clark | Lotus | 1990 | 1991 |

Example: Find the employees who worked during all the time that Brown worked.

One difficulty in expressing the above query in standard SQL is that the intervals are broken up into smaller ones. In the above Employee relation, Anderson worked at AT&T from 1980 to 1997, which is a single interval, but it is broken up into two separate records with intervals 1980 to 1993 and 1994 to 1997.

Querying TQuel Temporal Databases

coalesce the temporal intervals in the Employee relation.

$coalesce(2^{[I^-, I^+]}) \rightarrow 2^{[J^-, J^+]}$: This operator takes as input a set of intervals and returns another logically equivalent set of intervals, such that in the returned set no interval meets or overlaps another interval.

```
CREATE VIEW    Coalesce_Employee(Name, Company, From, To)
SELECT        Name, Company, Coalesce(From, To)
FROM          Employee
GROUP BY      Name, Company
```

Example: Find the employees who worked during all the time that Brown worked.

```
CREATE VIEW    Contemporary_of_Brown(Name)
SELECT        DISTINCT E1.Name
FROM          Coalesce_Employee AS E1, Coalesce_Employee AS E2
WHERE        E1.Name  $\neq$  "Brown" AND
            E2.Name = "Brown" AND
            contains(E1.From, E1.To, E2.From, E2.To)
```

Note the use of *contains* in the last line.

- TQuel uses SQL extended with the *coalesce* operator and *Allen's relations*.

Practice

1. Andrew and Barbara are two persons who have some free days next month as shown in the constraint relation *Free*, where d is an integer representing days of the month.

Free

| PERSON | DAY | |
|---------|-----|------------------------|
| Andrew | d | $3 \leq d, d \leq 6$ |
| Andrew | d | $16 \leq d, d \leq 26$ |
| Barbara | d | $6 \leq d, d \leq 18$ |
| Barbara | d | $25 \leq d, d \leq 30$ |

- (a) Convert the *Free* relation to the TQuel temporal data model.
 - (b) Convert the *Free* relation to the relational data model.
2. Answer the following queries using the *Free* relation converted to the TQuel temporal data model and interval-based temporal database queries.
 - (a) Find the days when Andrew and Barbara are both free.
 - (b) A common friend is visiting town from days 5 to 15. Find whether either Andrew or Barbara can be with the common friend on each of the days of the visit.
 3. Answer the queries in the previous exercise using the *Free* relation converted to the relational data model.